

# Munster Programming Training

## Cycle 2

## Lecture 2

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MPT 2-3-4 2016-17			
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November

Remember to change it

# Last lecture

- We introduce what an algorithm is
- We saw how a problem can be solved with different algorithms
- Not all the algorithms are equivalent

# Recall on Big O

Big O notation is used in Computer Science to describe the performance or complexity of an algorithm.

Big O specifically describes the worst-case scenario, and can be used to describe the execution time required or the space used (e.g. in memory or on disk) by an algorithm.

It is always a function of the problem size.

# Problem size $n$

Usually we call the problem size  $n$ .

$n$  can be:

- length of the array
- number of edge in a graph
- number of product to produce

In different occasions  $n$  can assume different meaning.

# Constant – $O(1)$

$O(1)$  describes an algorithm that will always execute in the same time (or space) regardless of the size of the input data set.

Example:

```
>>> 5 + 3
```

```
8
```

```
>>> k = k + 1
```

# Logarithmic – $O(\log N)$

It's early to understand  $O(\log N)$ . But consider that is slower than  $O(1)$  and faster than  $O(N)$ .

# Linear – $O(N)$

$O(N)$  describes an algorithm whose performance will grow linearly and in direct proportion to the size of the input data set.

```
>>>for x in range(0, n + 1):  
>>>    print "We're on time %d" % (x)
```



# Quadratic— $O(N^2)$

$O(N^2)$  represents an algorithm whose performance is directly proportional to the square of the size of the input data set.

This is common with algorithms that involve nested iterations over the data set.

```
>>>for x in range(1, n):  
>>>    for y in range(1, n):  
>>>        print '%d * %d = %d' % (x, y, x*y)
```

# Exponential– $O(2^N)$

$O(2^N)$  denotes an algorithm whose growth doubles with each addition to the input data set.

The growth curve of an  $O(2^N)$  function is exponential - starting off very shallow, then rising meteorically.

# Example

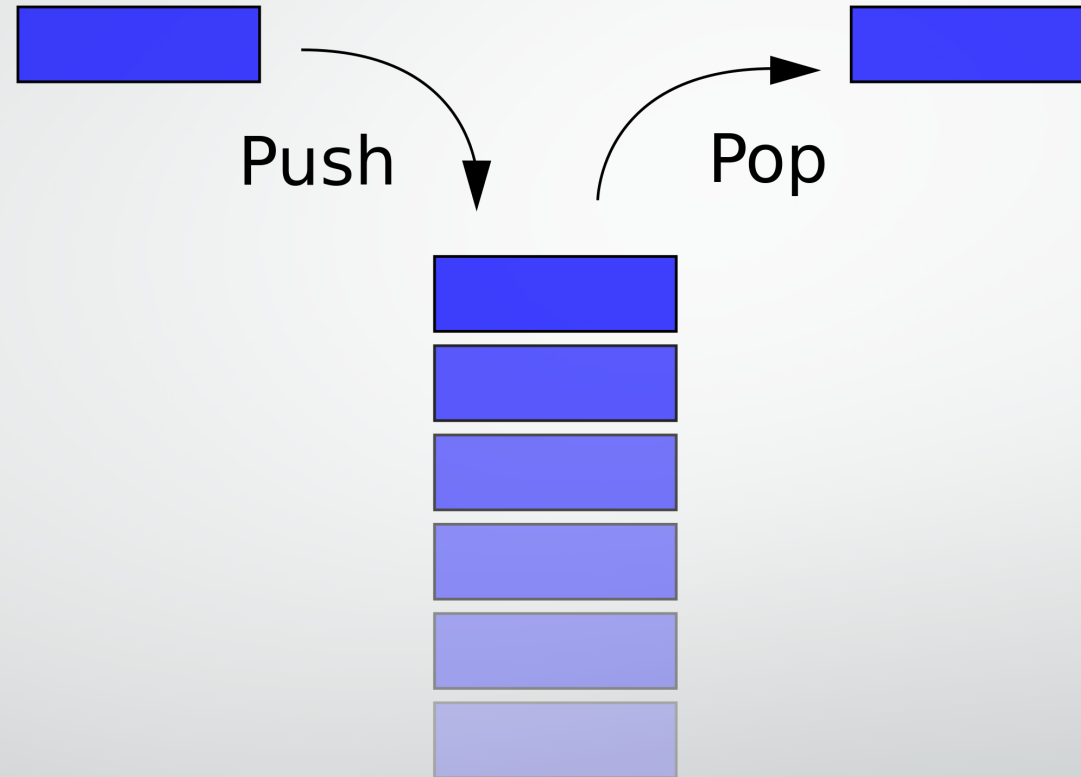
<http://bigocheatsheet.com/>



# Data Structure

It's important to choose the right one for your problem.

# Queue

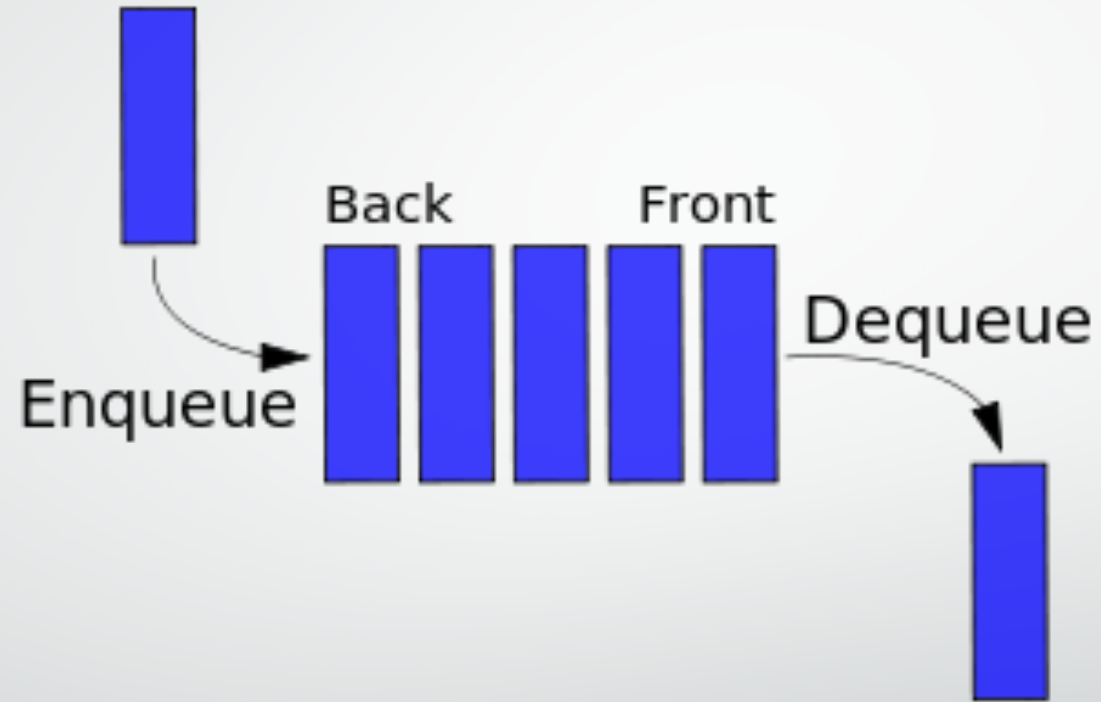


Order: Last in, first out (LIFO)

# Stack Example

```
>>> myStack = []  
>>>  
>>> myStack.append(3)  
>>> myStack.append(5)  
>>>  
>>> myStack  
[3, 5]  
>>> myStack.pop()  
5  
>>> myStack  
[3]
```

# Queue



Order: First in, first out (LIFO)

# Stack Example

```
>>> from collections import deque
```

```
>>> myQueue = deque([1, 2, 3])
```

```
>>> myQueue.append(4)
```

```
>>> myQueue.append(5)
```

```
>>> myQueue.popleft()
```

```
1
```

```
>>> myQueue.popleft()
```

```
2
```

```
>>> myQueue
```

```
deque([3, 4, 5])
```



# Exercise Stack and Queue

Your task in this exercise is to implement both the stack and queue operations described in the lecture. Your program should read in a list of integers and insert each one into both a stack and a queue. It should then print out the contents of the stack and then the contents of the queue.

## Input

The first line of input contains  $N$ , the number of integers.  $0 < N < 10000$ . The next line of input then contains  $N$  integers.

## Output

Your program should output the contents of the stack in last-in first-out order, followed immediately by a new line. Then the contents of the queue in first-in first-out order.